

# Foundation & Challenges of Quantum Computing

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### Random Bits



Classical randomized algorithms use classical coins (random bits).

### **Classical randomized algorithms are efficient**

Quicksort, Monte-Carlo simulation, sampling, ...

Satisfiability of 3SAT Boolean formula

Brute force O(2<sup>n</sup>)

Deterministic O(1.439<sup>n</sup>)

Randomized O(1.321<sup>n</sup>)

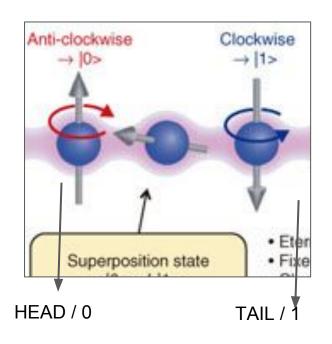
Remote verification of equality of bit strings

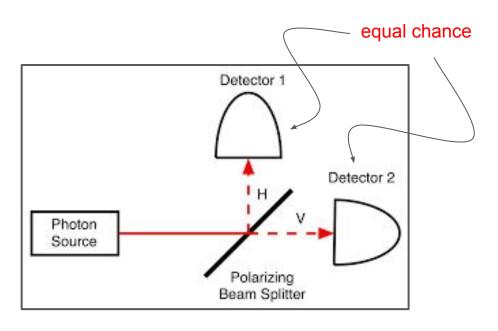
Brute force O(n)

Randomized O(log n)

### **Qubits**

Photon polarization, Electron spin, direction of current in Joesphon junction, ...



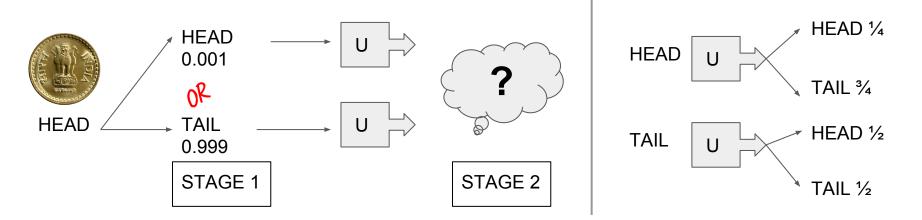


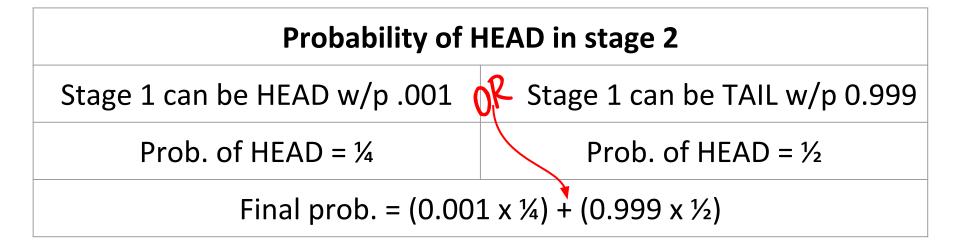
Quantum operations are inherently probabilistic.

Quantum ≈ Randomized +++++

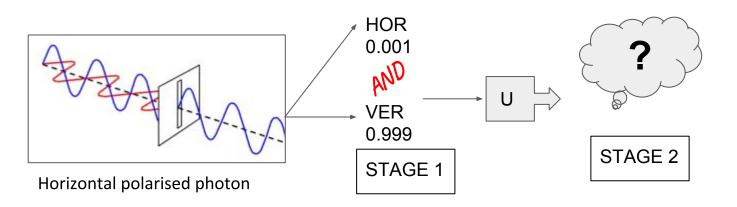
(different rules of randomization)

### Probabilities of Classical states

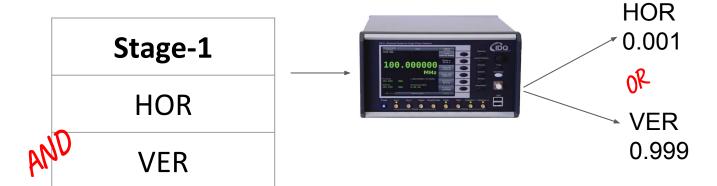




## Probabilities of Quantum states



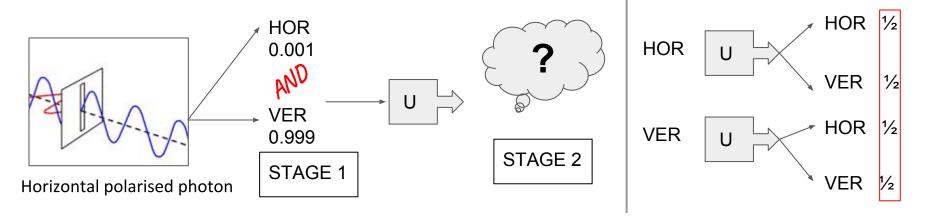
superposition

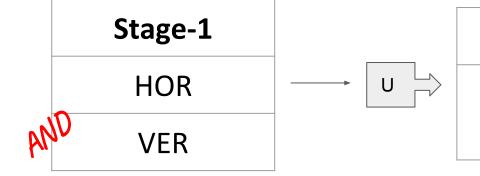


collapse under measurement

### Probabilities of Quantum states

Observation probabilities

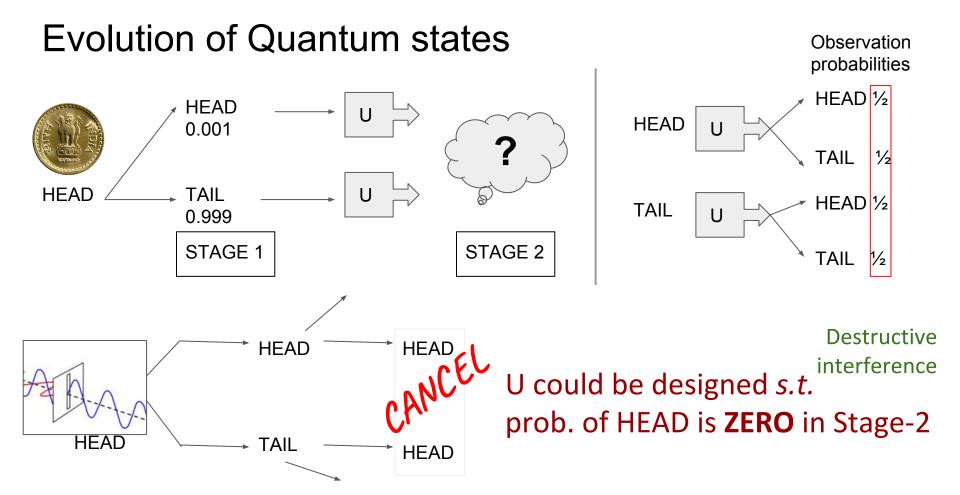




### **Probability of HOR in stage 2**

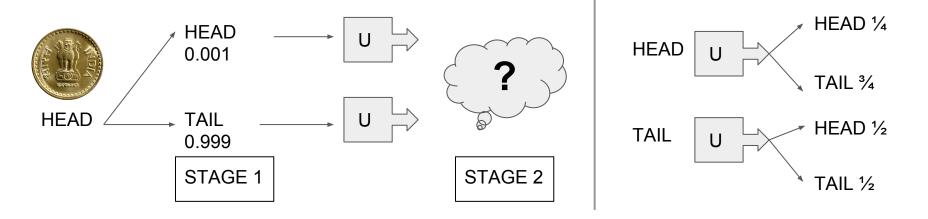
Prob. of HOR from H in stage-1

+ Prob. of HOR from V in stage-1



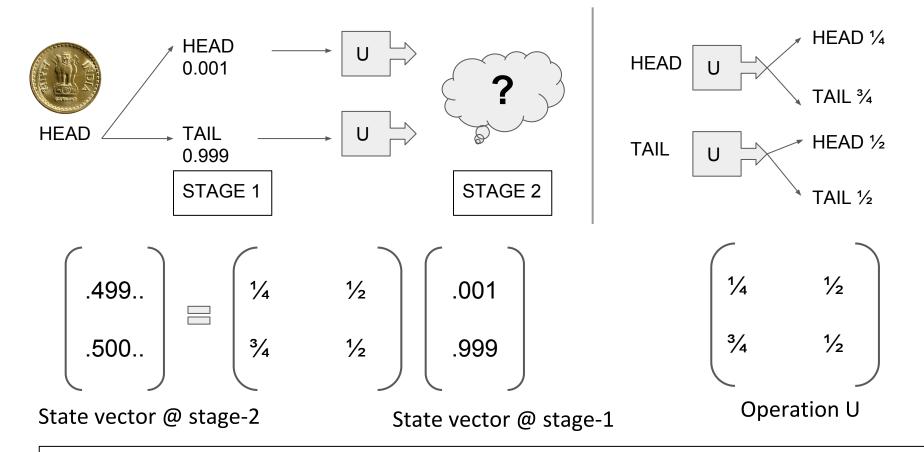
Classical: State is a unit L1-norm real vector, U is stochastic Quantum: State is a unit L2-norm complex vector, U is unitary

### **Evolution of Classical states**



$$\begin{pmatrix}
0.001 \\
0.999
\end{pmatrix}
= 0.001 \begin{pmatrix}
1 \\
0
\end{pmatrix}
+ 0.999 \begin{pmatrix}
0 \\
1
\end{pmatrix}$$
State vector @ stage-1 HEAD TAIL

### **Evolution of states**



Classical: State is a unit L1-norm real vector, U is stochastic

Quantum: State is a unit L2-norm complex vector, U is unitary

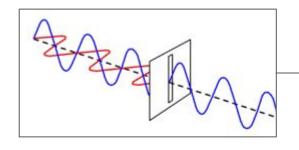
# Quantum Searching in Database

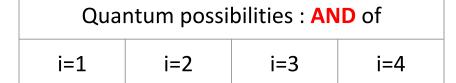
A[1]	A[2]	A[3]	A[4]
0	0	1	0

**Task:** Find some x s.t. A[x] = 1.

Any A[j] can be queried.

Efficiency measure : no. of queries.





Query A[i]

Always ...

A[3]=1

Other possibilities cancel each-other

Grover's amplitude amplification

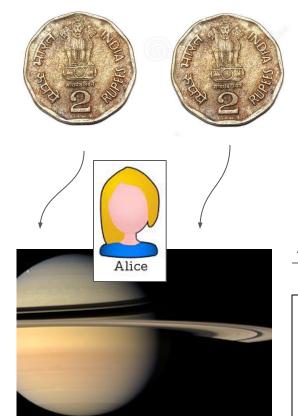
Quantum possibilities : AND of

A[1]=0 A[2]=0 A[3]=1 A[4]=0

Measure

Get A[3]=1 with prob.  $\frac{1}{4}$ 

# Secret Key Distribution

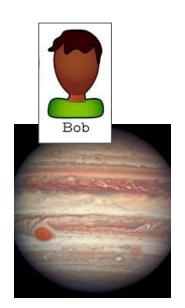


Two coins.

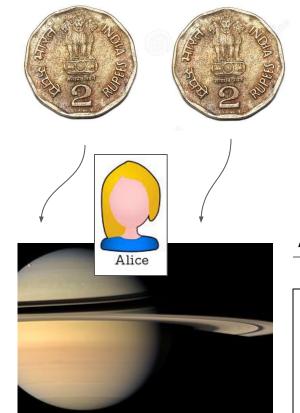
HH or TT with equal probability.

Alice throws one coin to Bob

- 1. Alice checks her coin.
- 2. Bob checks his coin.
- Alice and Bob share a random bit.
- 4. This can be repeated for multi-bit key.



# Secret Key Distribution (with Eve)



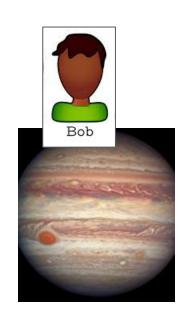
Two coins.

HH or TT with equal probability.

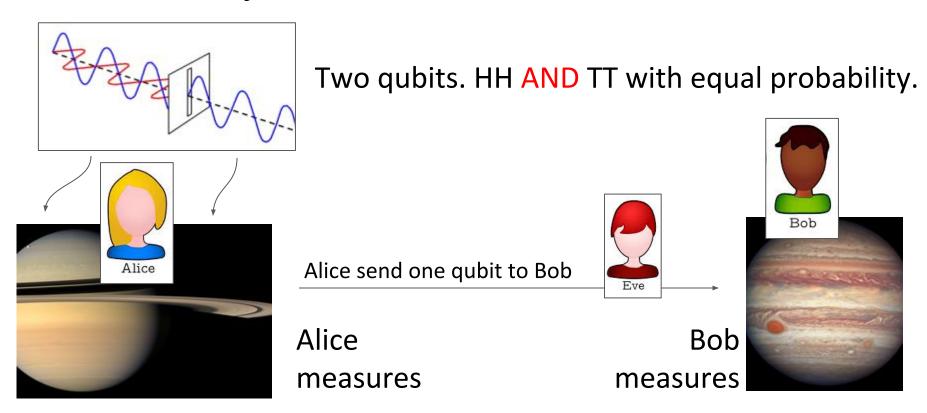




- 1. Alice checks her coin.
- 2. Bob checks his coin.
- Alice and Bob share a random bit.
- 4. Eve also knows the shared bit.



# Quantum Key Distribution



If Eve intercepts and measures, say H, Alice and Bob get HH.

If Eve does not intercept, Alice & Bob gets HH or TT with prob. ½.

Alice & Bob run random verification tests to detect Eve.

Possible since state of qubit "collapses" during measurement.

# Where are the quantum Browsers and Softwares?

Qubits cannot be copied!

```
is the code that does the bubble sort.
for (int i = ar.length - 1; i > 0; i--) {
   for (int j = 0; j < i; j++) {
      if (ar[j] > ar[j + 1]) {
      temp = ar[j];
      ar[j] = ar[j + 1];
      ar[j + 1] = temp;
```

Quantum variables hold multiple values at the same time!

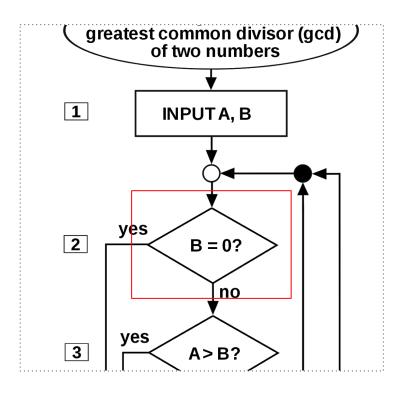
**Output**: A quantum state proportional to  $|x\rangle$  where  $x \approx x^- = A$  **Algorithm**:

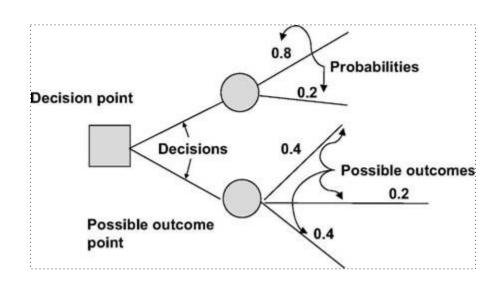
- 1. Prepare the quantum state  $|b\rangle = \frac{1}{\|b\|_2} \sum_{i=1}^{n} \beta_i |v_i\rangle$ .
- 2. Perform phase estimation to create the state  $\frac{1}{1000}\sum_{i=1}^{n} \beta_i | 1$

# Where are the quantum Browsers and Softwares?

Measuring qubit changes its state.

Final probability is NOT the sum of sub-tree probabilities.





# **Promises and Prospects**

**1981** Feynman proposed quantum computer to efficiently simulate many-body quantum systems

- 1984 Bennett and Brassard designed quantum protocol for BB84 running on 2000KM secret key sharing
- fiber-optic cable in China 1991 Another QKD protocol by Ekert

QKD networks: DARPA, Tokyo, Vienna, Japan, ...















# **Promises and Prospects**

Technology not clear

**1985** Deutsch proposed a general purpose programmable quantum computer

Quantum
supremacy
race

**1992** Deutsch and Jozsa solved a (toy) problem in half the time taken by the best classical algorithm

1993 Simon designed algorithm that is efficient on quantum computer but inefficient classically

1994 N is the best

**1996** Grover designed algorithm to search in a database of N elements using  $\sqrt{N}$  "lookups" (classical best is N/2)

... better-than-classical algorithms for problems on numbers, graphs, geometric objects, strings, statistics, communication, data structures, ... but <u>limited</u> speedup

# **Promises and Prospects**

2001 15 factored using 10<sup>18</sup> identical molecules

- Requires 1994 Shor designed algorithms to factor n-bit number in high-precision O(n<sup>2</sup>) time (classical best is O(exp(n<sup>1/3</sup>)))
  - 1995 Shor and Steane designed error-correcting codes
  - 1998 Gottesman and Knill showed how to efficiently Oops! simulate certain quantum algorithms classically
- 2017 Microsoft releases 40-qubit classical simulator
- ODE, PDE, 2009 Harrow+ designed linear system "solver" Quantum ML
- 2015 Grassl showed 3000-7000 qubits needed to search cryptography AES key using Grover's algorithm
  - 2016 Google simulated a Hydrogen molecule with 9 qubits

# Summary

Quantum mechanics that drive quantum computing is mysterious

But if you are a believer ... (\*)

Quantum algorithm design and analysis possible using knowledge of algorithms, probability and linear algebra.

Thanks to physicists, material scientists, engineers, mathematicians, ... in universities, R&D labs and corporates ...

These algorithms can be implemented on real quantum computers and experimented with.

Too early to say how and where QC will become useful ...

Just the right time to enter the game.

### Gartner Hype Cycle for Emerg



Tin

### gartner.com/SmarterWithGartner

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